

# Unbalanced PSI

## Applications, Constructions, and Combinations with PIR

Dr Christian Weinert – Presented at WPEC 2024, September 24  
NIST Workshop on Privacy-Enhancing Cryptography



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# The Cryptography Group



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- RHUL has a long history of cryptographic research, dating back to at least **1987**
- Currently:
  - 11 members of staff
  - ~15 PhD students



- Find out more: <https://cryptography.isg.rhul.ac.uk/people/>

# Definition

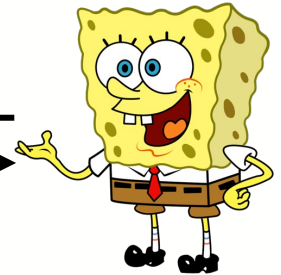
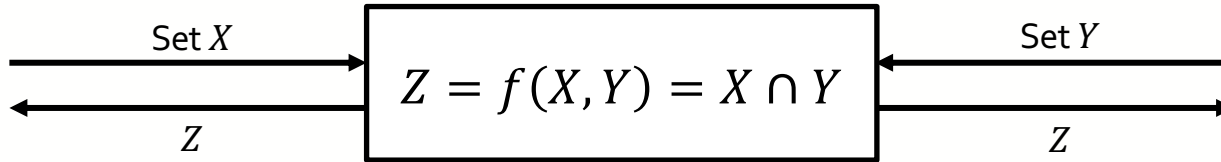


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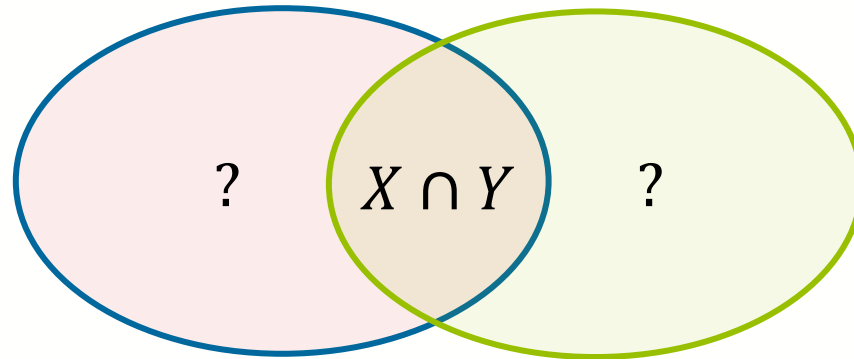
# Private Set Intersection (PSI)



Alice



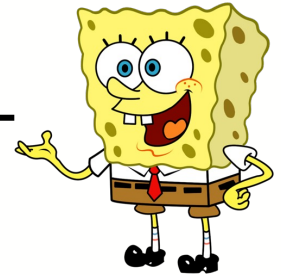
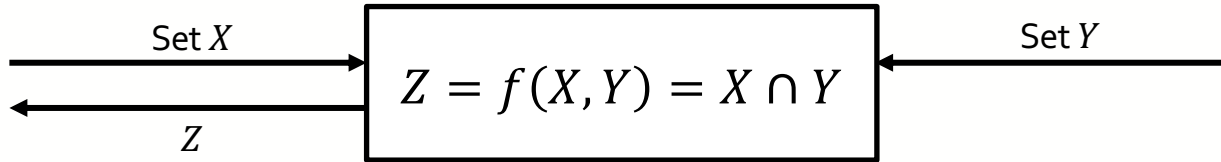
Bob



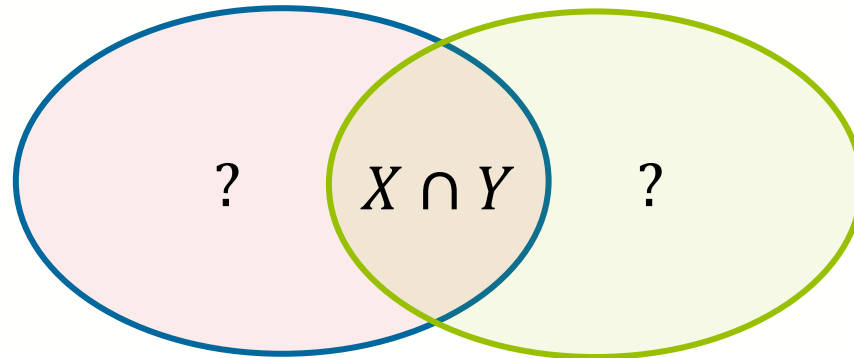
# Private Set Intersection (PSI)



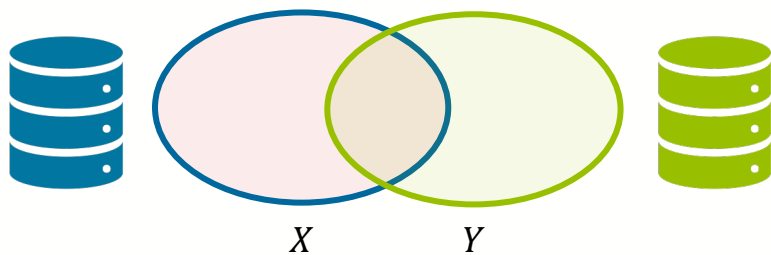
PSI Receiver



PSI Sender

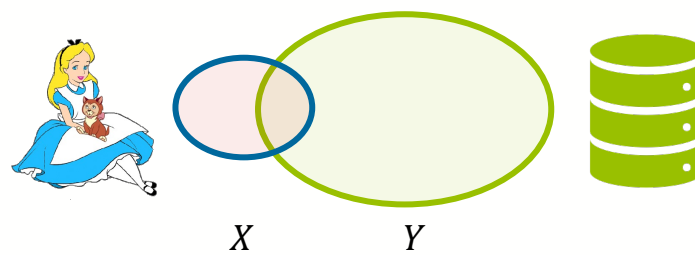


# Balanced vs Unbalanced PSI (uPSI)



Balanced PSI

$$\begin{array}{c} \uparrow \\ |X| \approx |Y| \end{array}$$



Unbalanced PSI (uPSI)

$$\begin{array}{c} \uparrow \\ |X| \ll |Y| \end{array}$$

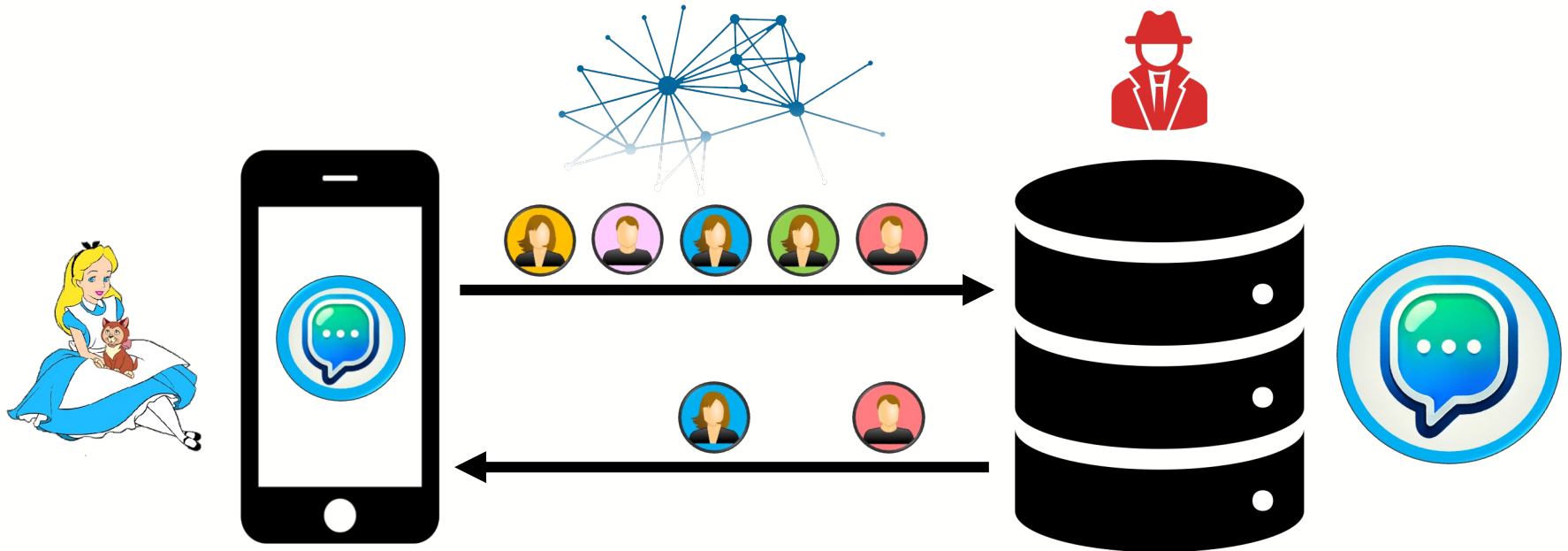
# Applications and Deployments



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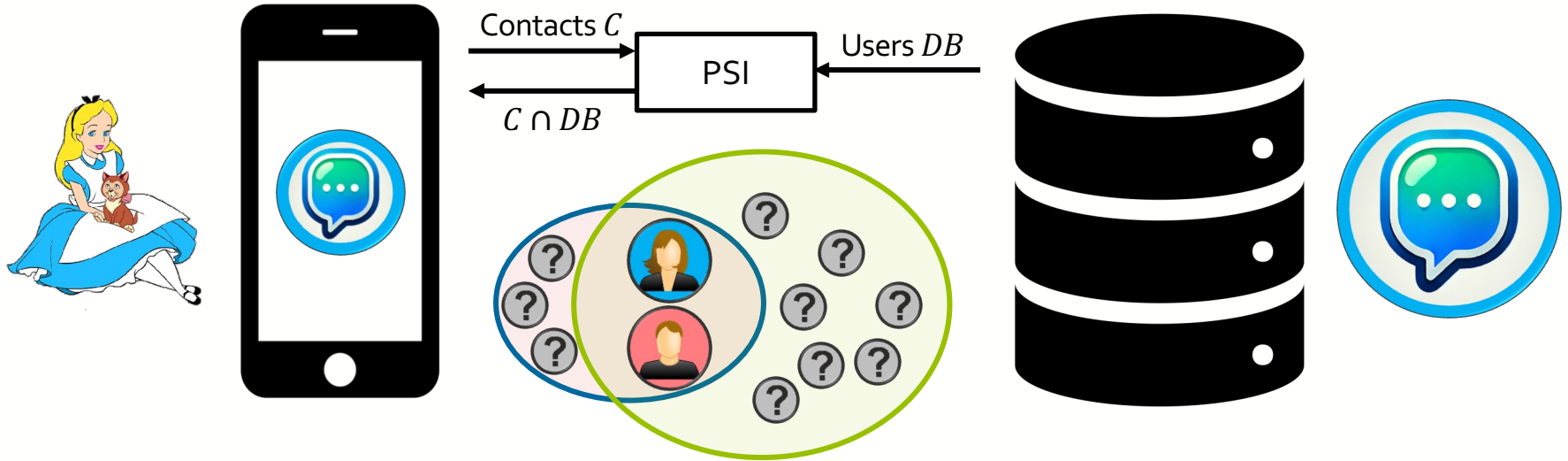


# Mobile Contact Discovery



<https://contact-discovery.github.io/>

# Mobile Contact Discovery



<https://contact-discovery.github.io/>

# Intel SGX-based Contact Discovery Service in Signal

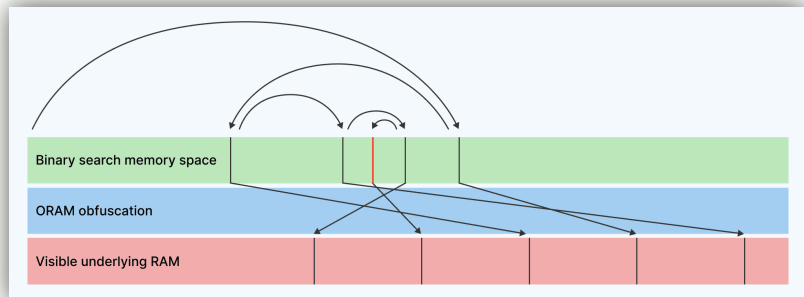


Performance requirements: 1B+ users, 10k contacts, <2s latency, <10MiB communication

## CDSI SGX Enclave

The Contact Discovery Service (CDSI) SGX Enclave provides endpoints to expose a table of Signal user records indexed by phone number that allows users to securely discover which of their contacts are also Signal users.

While SGX provides memory encryption and attestation that are essential to the security of this enclave, it does not guard against memory access pattern or timing side-channels. To close memory-access pattern side channels at the architectural level, the table of user records is stored in [Oblivious RAM \(ORAM\)](#). To close timing and access-pattern side-channels at the local level, the code is written using branchless idioms and oblivious algorithms. Please see [Side-channel Resistant Programming Idioms](#) before contributing.



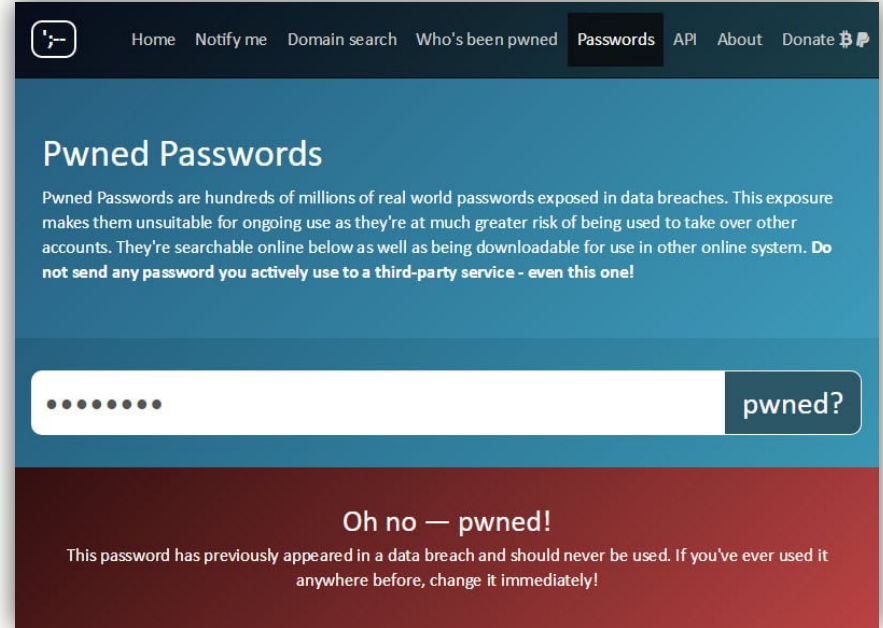
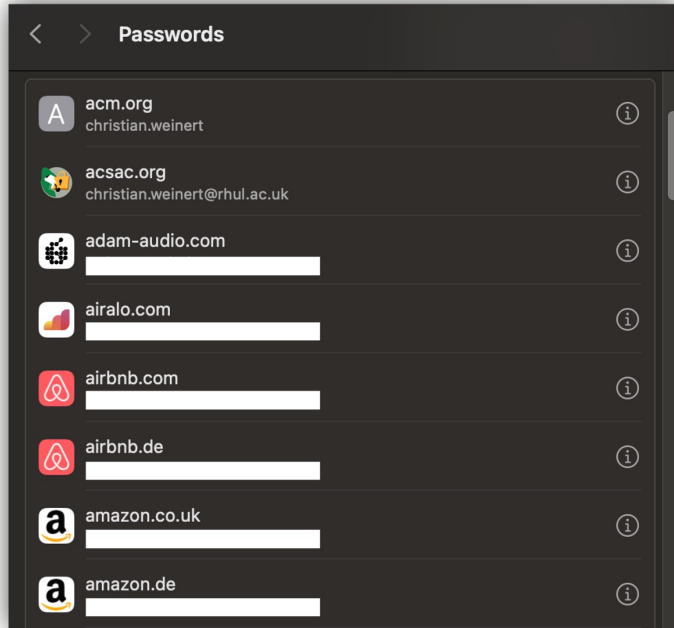
Path ORAM [Sig22]

**Mark Ermolov**  
@markel\_

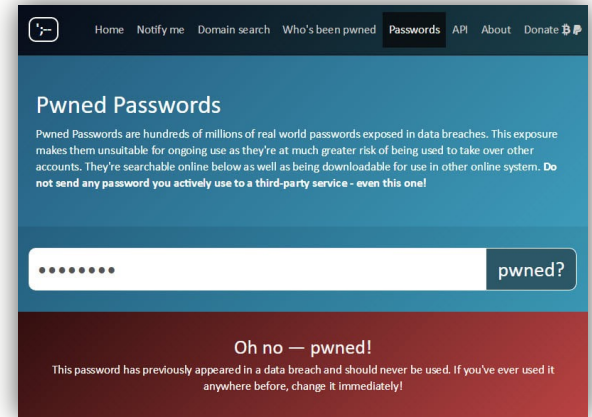
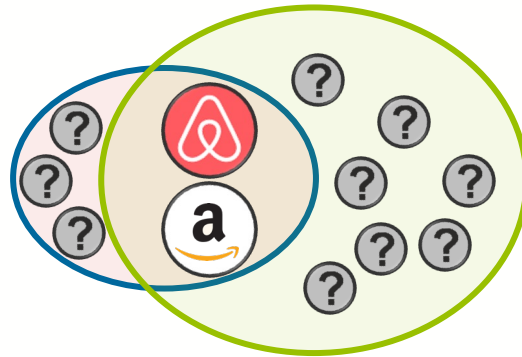
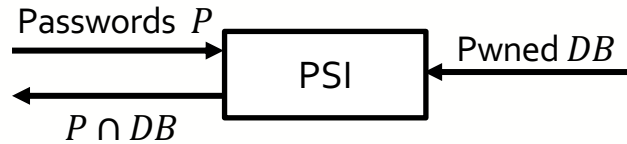
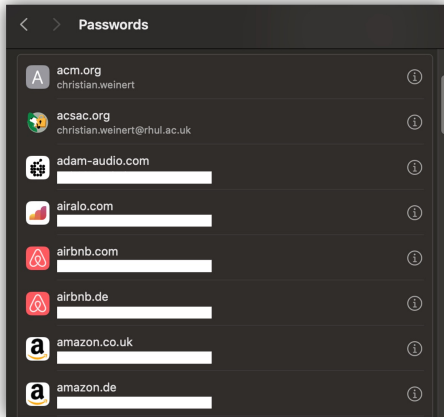
Intel HW is too complex to be absolutely secure! After years of research we finally extracted Intel SGX Fuse Key0, AKA Root Provisioning Key. Together with FK1 or Root Sealing Key (also compromised), it represents Root of Trust for SGX. Here's the key from a genuine Intel CPU 😊

```
Administrator: Command Prompt - python
>>> ipc.indrscan(dfx_agg, 0x51, 64)
[64b] 0x00131A0000007443
>>> ipc.threads[0].halt()
>>>
[GLP_C0_T0] Halt Command Break at [0x38:00000000744f0193] -- 18:53:46.922472 2024-08-26
[GLP_C1_T0] Halt Command Break at [0x38:000000000009e1de] -- 18:53:46.923484 2024-08-26
[GLP_C2_T0] Halt Command Break at [0x38:000000000009e1de] -- 18:53:46.924484 2024-08-26
[GLP_C3_T0] Halt Command Break at [0x38:000000000009e1de] -- 18:53:46.924484 2024-08-26
>>> ipc.threads[0].mem("0x0038484000", 4, 3)
>>> ipc.indrscan(dfx_agg, 0x51, 64)
[64b] 0x00131A0000004706C
>>> sb_print(0x1, 0xf000, 0xf000)
fc0b: 10713309 3300750e 40000000 0020067a
fc10: 00000000 00000000 00000000 0000ffff
fc20: 00c00000 00713704 00018037 0e760000
fc30: 30323a1a 00000000 00000023 00550770
fc40: 0e760a76 00000000 00000000 00000000
fc50: 00000000 00000000 7201c2c5 5c01d004
fc60: 2003e044 00000020 00000000 00000002
fc70: 83000000 10400000 00000000 00000002
fc80: 88000000 10400000 00000000 00000002
fc90: 88000000 10400000 44000002 02075fa5
fc0b: 0214600b 02c02004 00c2c000 2fa24400
fc0c: 608b0207 2094b214 000002c0 03070f00
fcc0: 3a162000 5326aab4 552aaa4a 1186514e
fc0b: 00000000 00000002 88000000 10400000
fc0c: 00000000 00000002 88000000 10400000
fcc0: c696e892 e27d1d13 cc860ff7 40b157a5
>>>
```

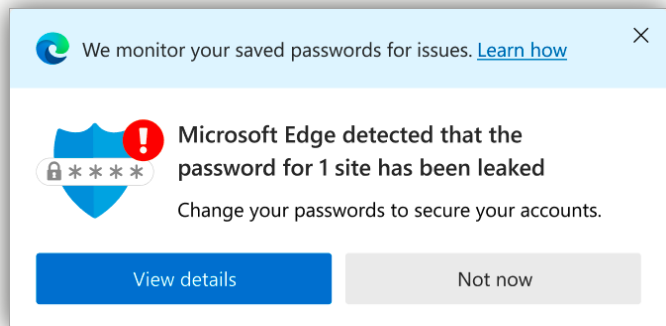
# Discovery of Leaked/Compromised Passwords



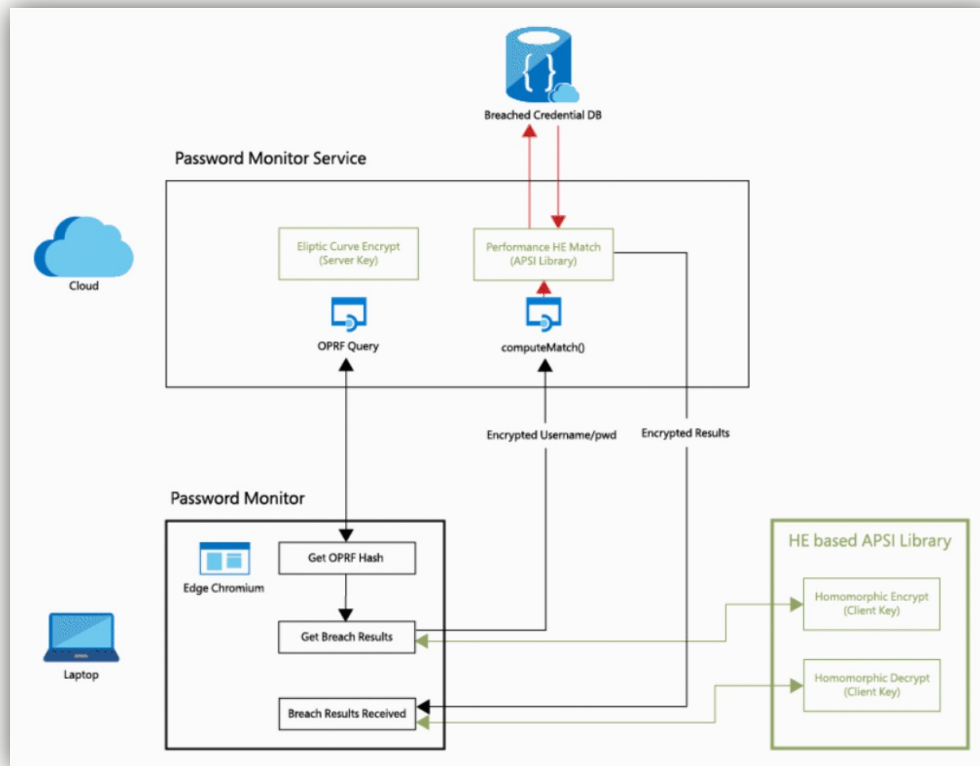
# Leaked/Compromised Passwords



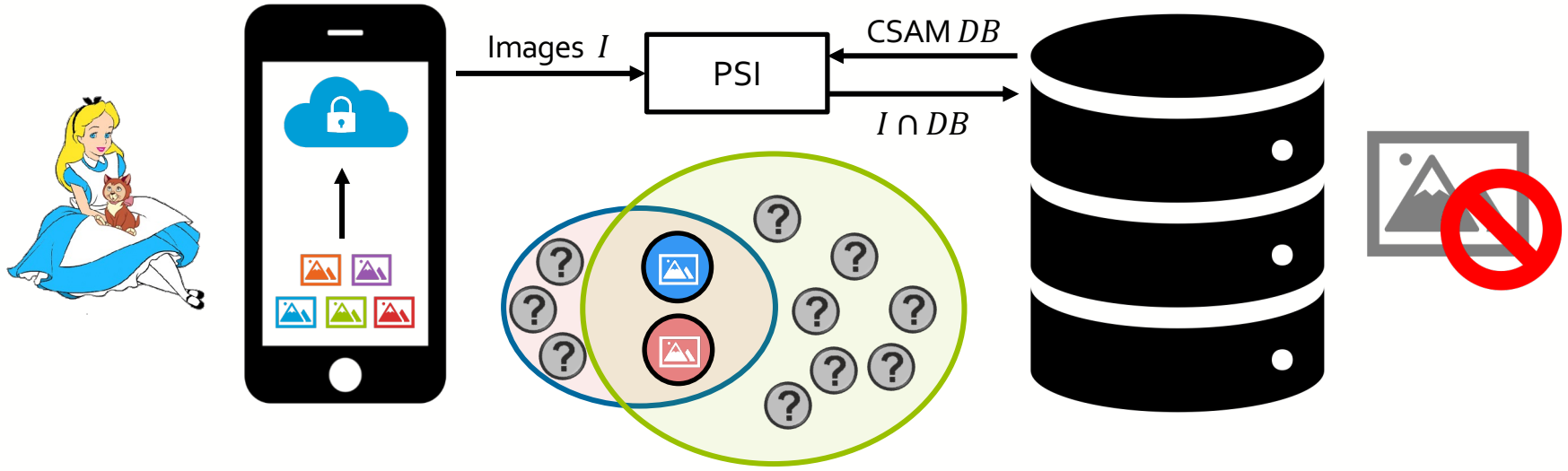
# Password Monitor in Microsoft Edge



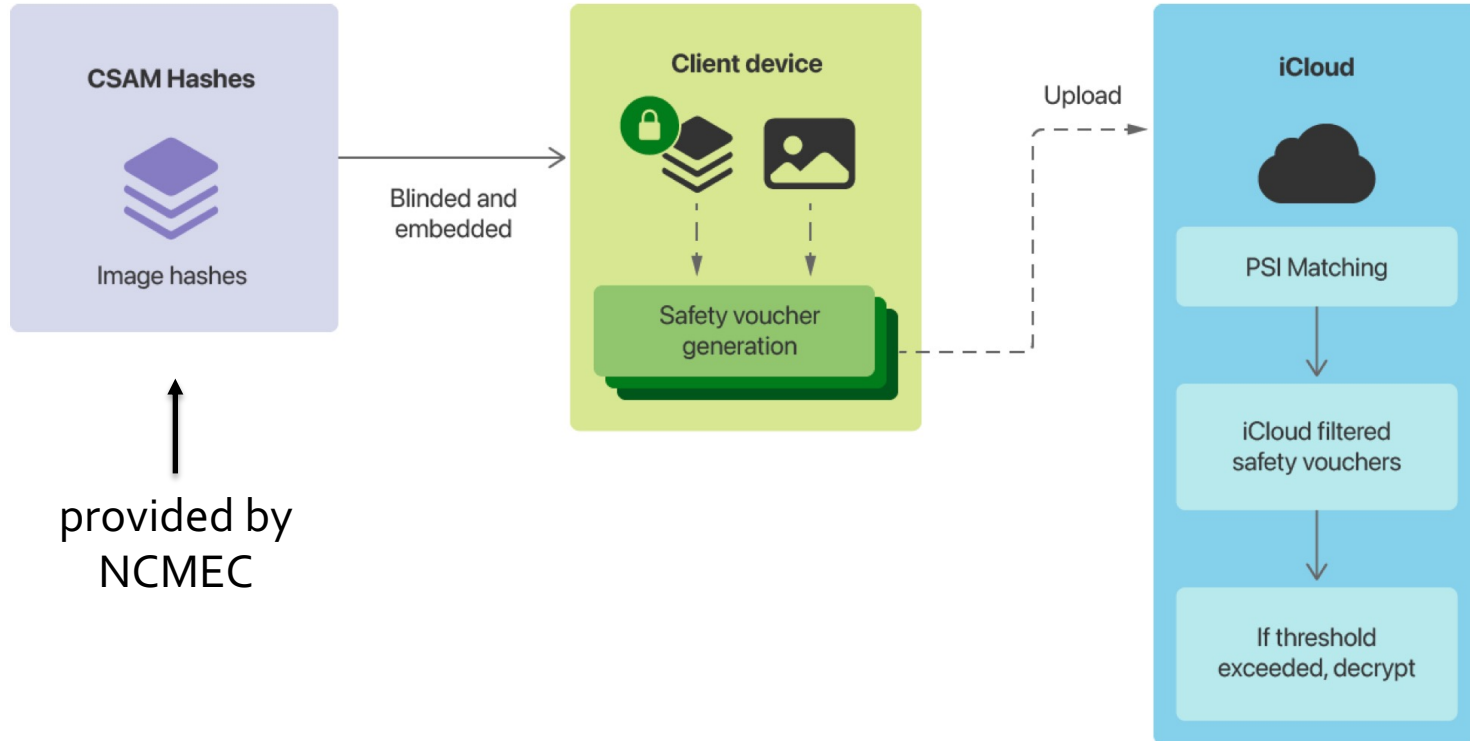
Framework based on  
FHE-based unbalanced PSI  
[Mic21]



# Scanning E2E-encrypted Content

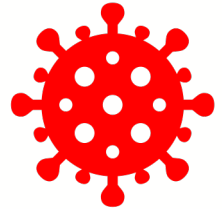
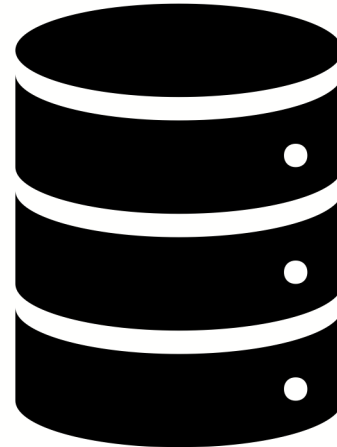
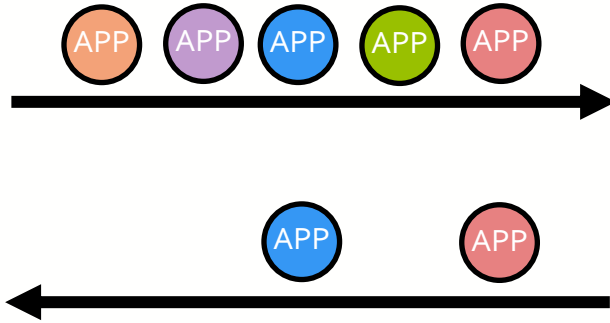
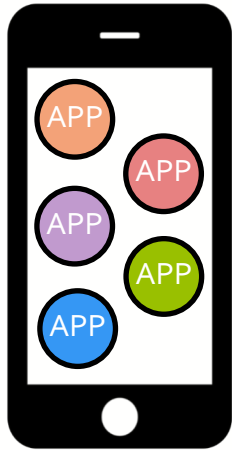


# Apple's iCloud CSAM Detection [App21]

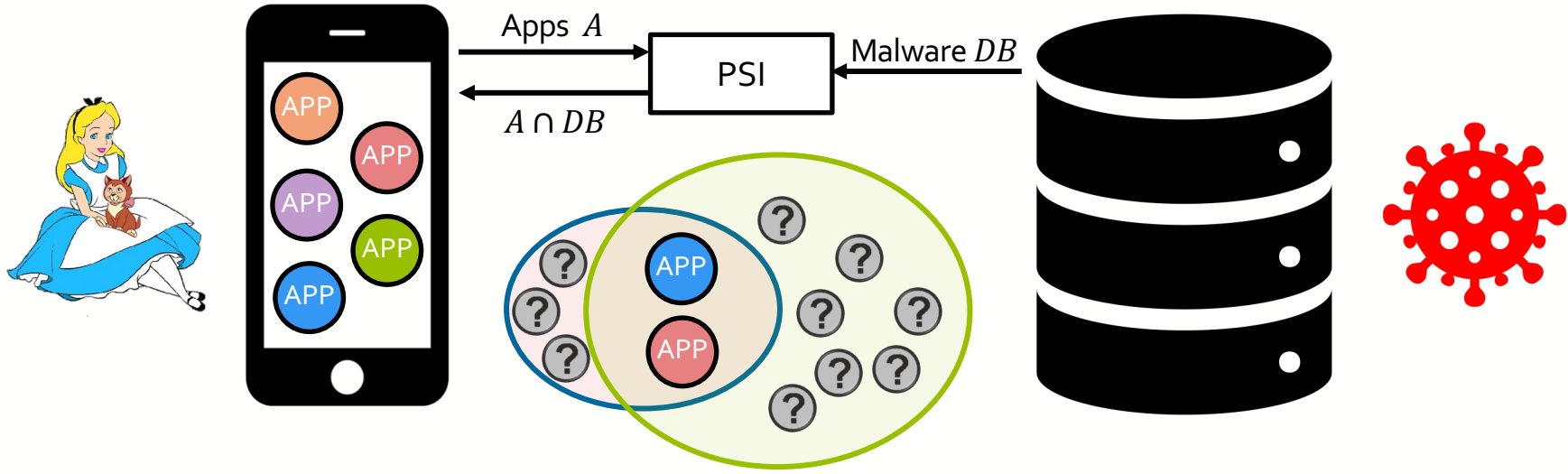




# Mobile Malware Detection Service [KLS+17]



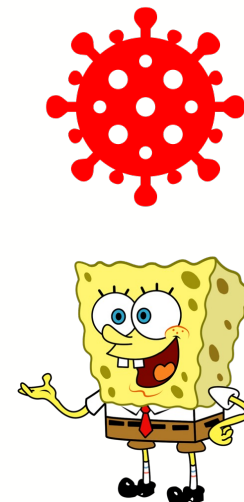
# Mobile Malware Detection Service [KLS+17]



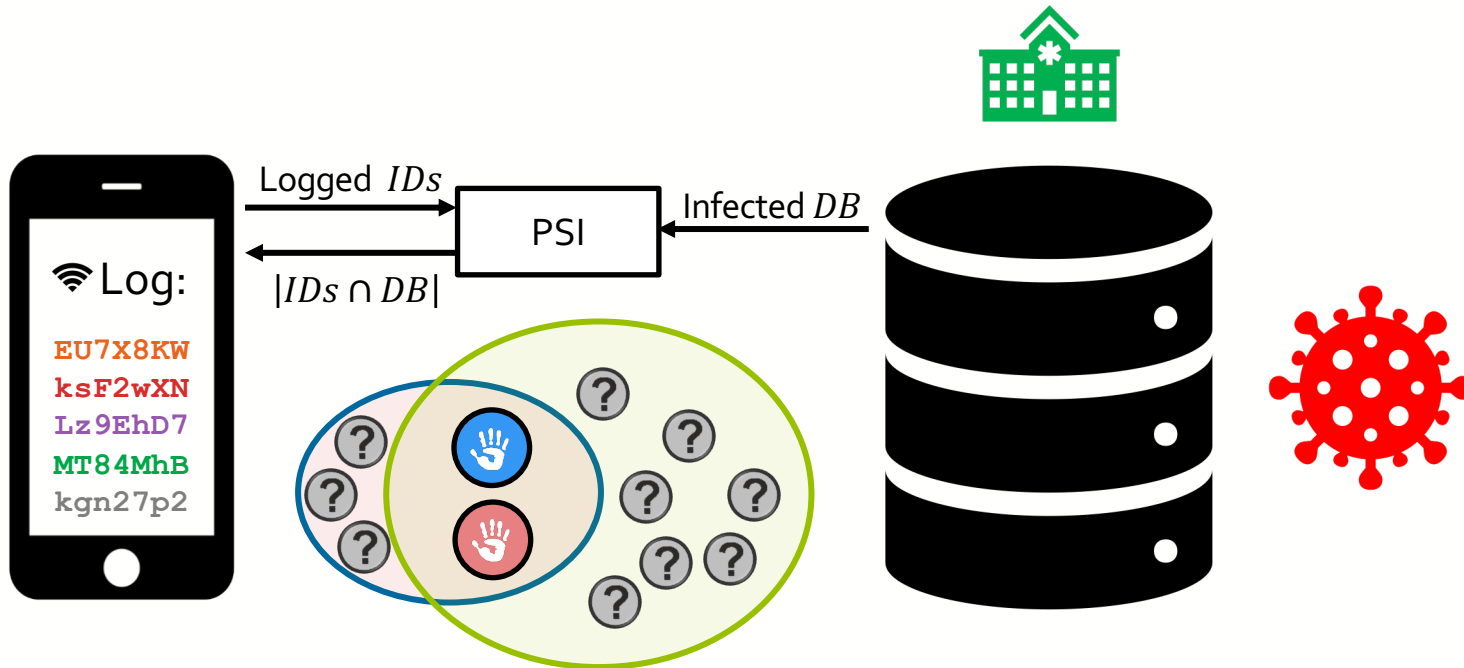
# Contact Tracing (via uPSI Variants)



Report EU7X8KW



# Contact Tracing (via uPSI Variants)



Relevant works: [DPT20, TSS+20, WY23]

# Constructions

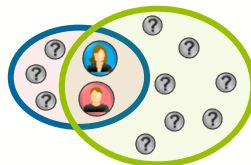


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# OPRF-based uPSI – Framework



Client inputs  $x_{i \in \{1, \dots, n\}}$



Precomputation Form suggested by [KLS+17]



## 1. Setup Phase

Client-independent Precomputation Phase,  $O(|Y|)$

Apply PRF to all server inputs with key  $k$  and insert them in probabilistic data structure  $PD$  for membership testing

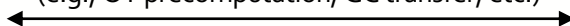
$PD$



## 2. Base Phase

Client-specific Setup Phase,  $O(|X|)$

(e.g., OT precomputation, GC transfer, etc.)



## 3. Online Phase ( $O(|X|)$ )

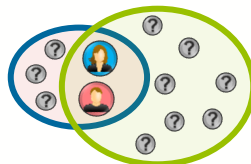
Check if  $e_i$  is in  $PD$



# OPRF-based uPSI – Instantiations



Client inputs  $x_{i \in \{1, \dots, n\}}$



## 1. Setup Phase

Client-independent Precomputation Phase,  $O(|Y|)$

Bloom Filter [KLS+17],  
Cuckoo Filter [RA18, KRS+19]

$PD$

Apply PRF to all server inputs with key  $k$  and insert them in probabilistic data structure  $PD$  for membership testing

## 2. Base Phase

Client-specific Setup Phase,  $O(|X|)$   
(e.g., OT precomputation, GC transfer, etc.)

OPRF Options  
[KLS+17, KRS+19, ...]:

- Naor-Reingold
- AES/LowMC GC
- RSA
- ...

## 3. Online Phase ( $O(|X|)$ )

$x_i$   
 $e_i$

OPRF

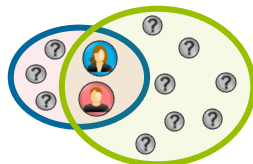
$k$

Check if  $e_i$  is in  $PD$

# OPRF-based uPSI – Performance for $|X| = 2^{10}$ , $|Y| = 2^{28}$



Client inputs  $x_{i \in \{1, \dots, n\}}$



## 1. Setup Phase

Client-independent Precomputation Phase,  $O(|Y|)$

[KRS+19]:  $\approx 1\text{GB} / 15\text{s}$

Apply PRF to all server inputs with key  $k$  and insert them in probabilistic data structure  $PD$  for membership testing

$PD$

## 2. Base Phase

Client-specific Setup Phase,  $O(|X|)$

(e.g., OT precomputation, GC transfer, etc.)

[KRS+19]:  $\approx 2\text{MB} / 0.1\text{s}$

## 3. Online Phase ( $O(|X|)$ )

[KRS+19]:  $\approx 2\text{MB} / 0.6\text{s}$

Check if  $e_i$  is in  $PD$

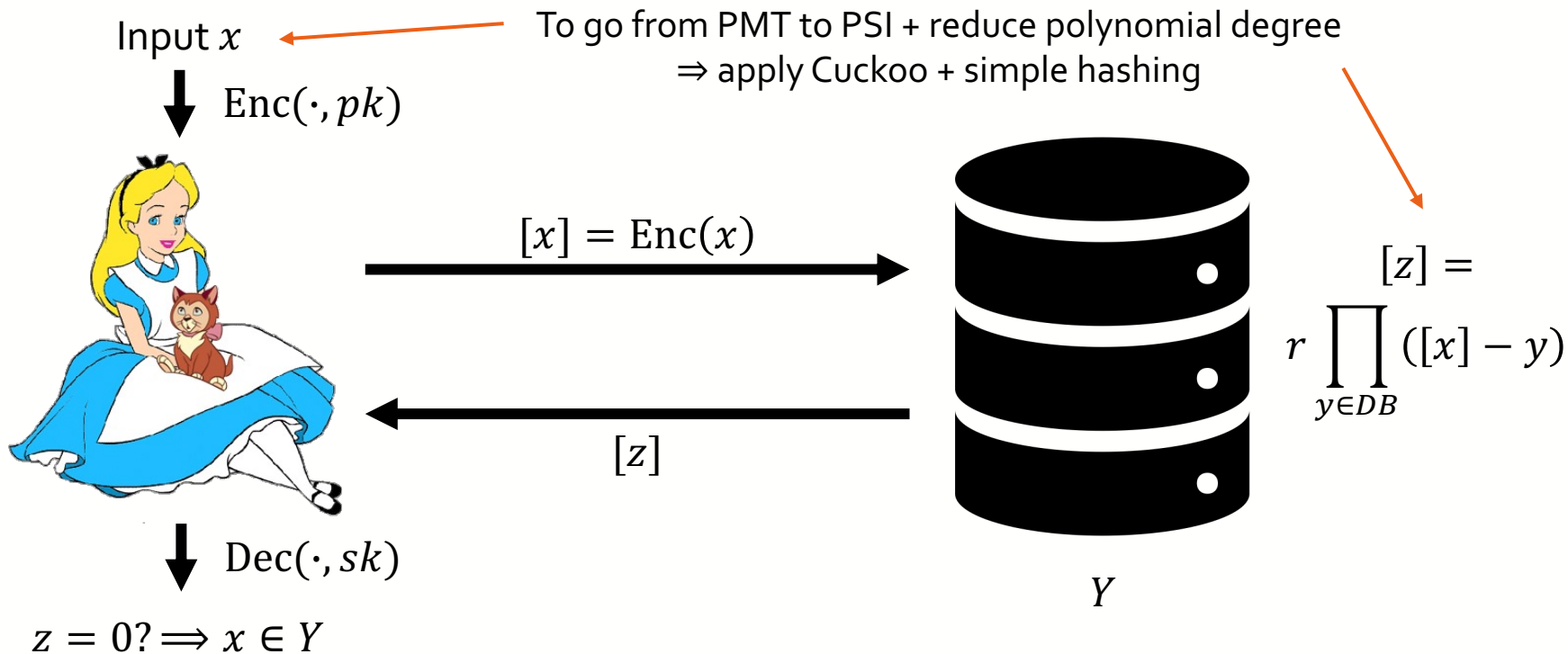
$x_i$   
 $e_i$

OPRF

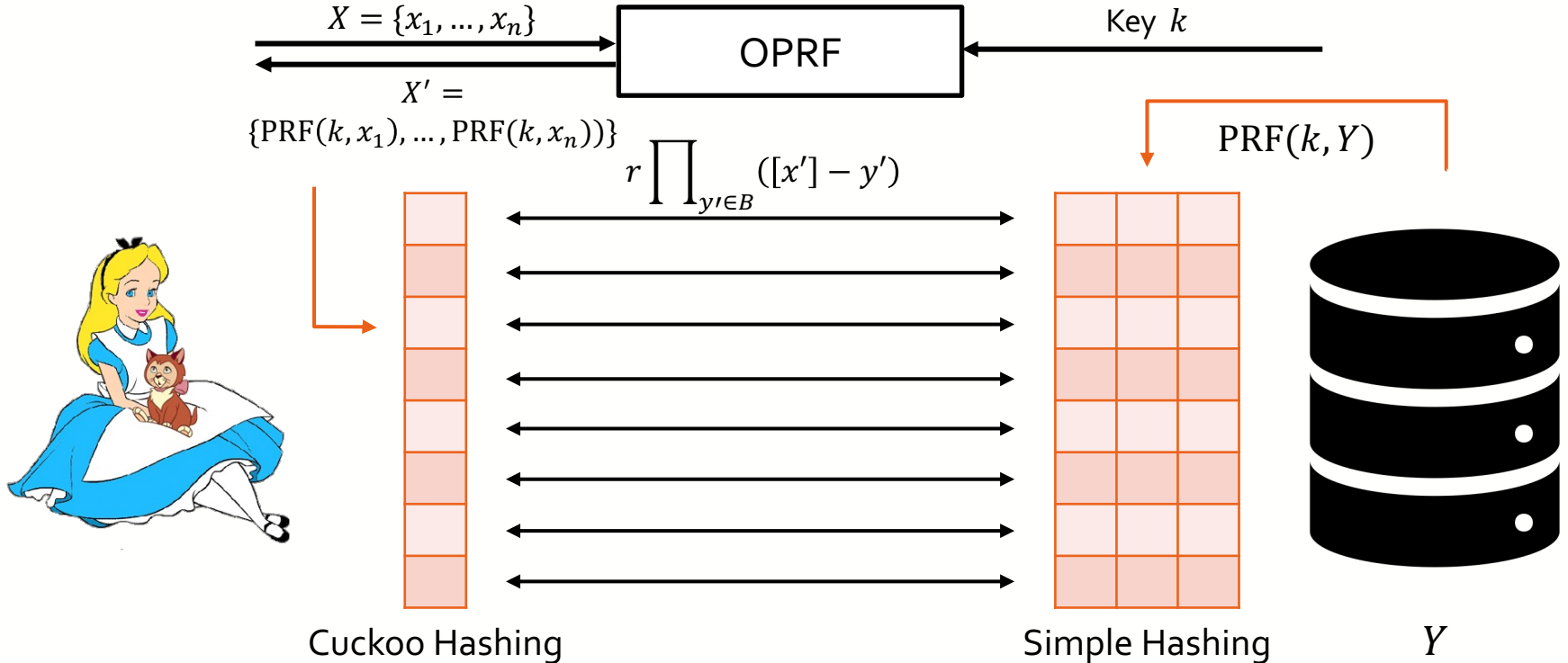
$k$

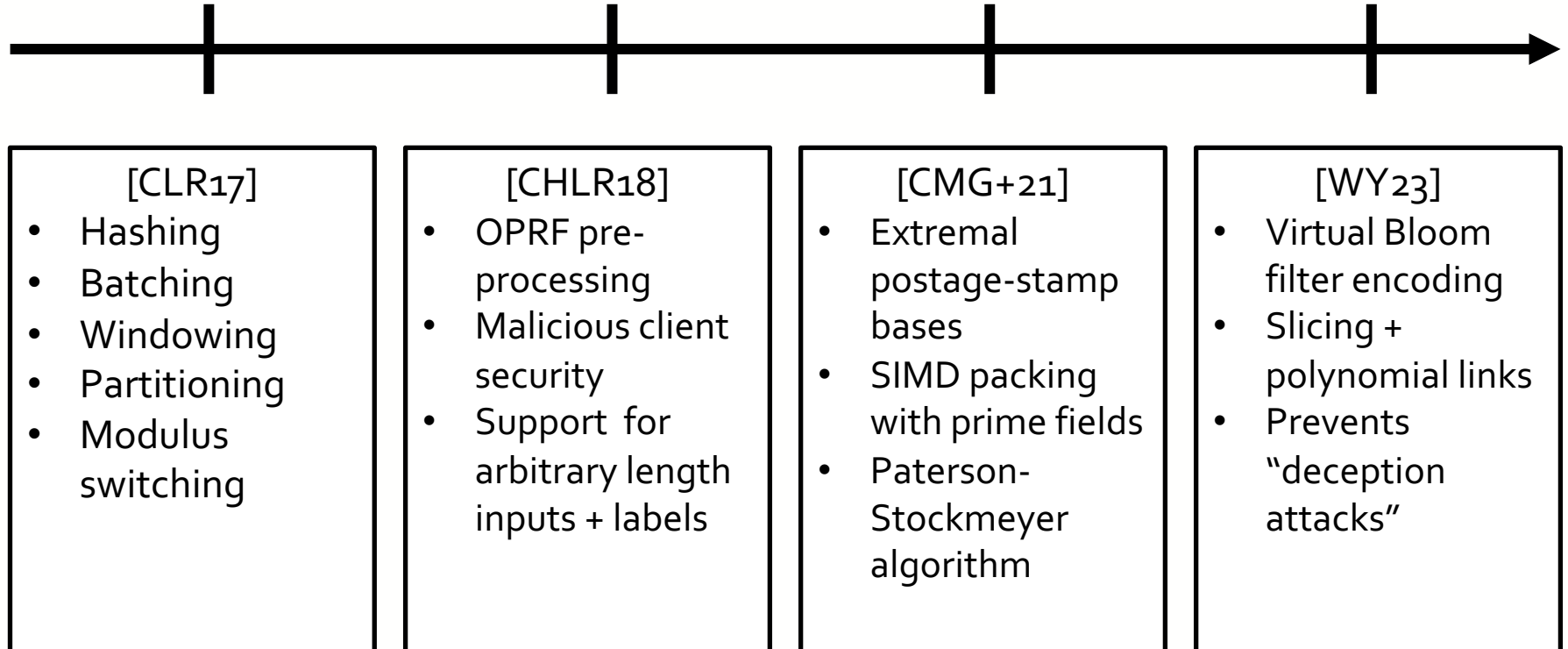


# FHE-based uPSI – General Idea



# FHE-based uPSI – Framework





# FHE-based uPSI – Performance for $|X| = 2^{10}$ , $|Y| = 2^{28}$



Protocol	Time [s]		Comm. [MB]
	Server Setup	Online	
[CHLR18]	4,628	12.1	18.57
[CMG+21]	2,033	8.21	12.56
[WY23]	3,298	8.00	5.22

All results are in LAN setting with 32 threads on the server side

Note: “PEPSI” [MLE+24] for circuit-based uPSI reduces computation complexity from  $O(|Y| \log|Y|)$  to  $O(|Y|)$  but is concretely less efficient for (labeled) uPSI (>100s online time)

# Combinations with PIR

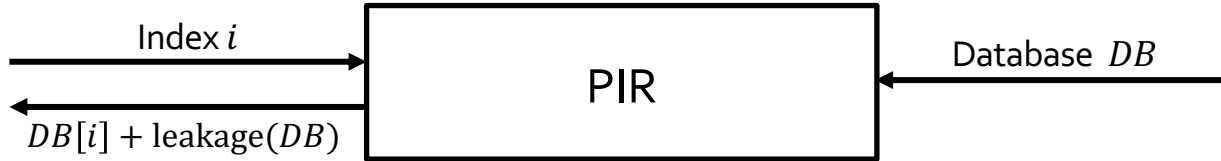


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# Private Information Retrieval (PIR)



Alice



Bob

non-trivial PIR: communication overhead  $< O(|DB|)$

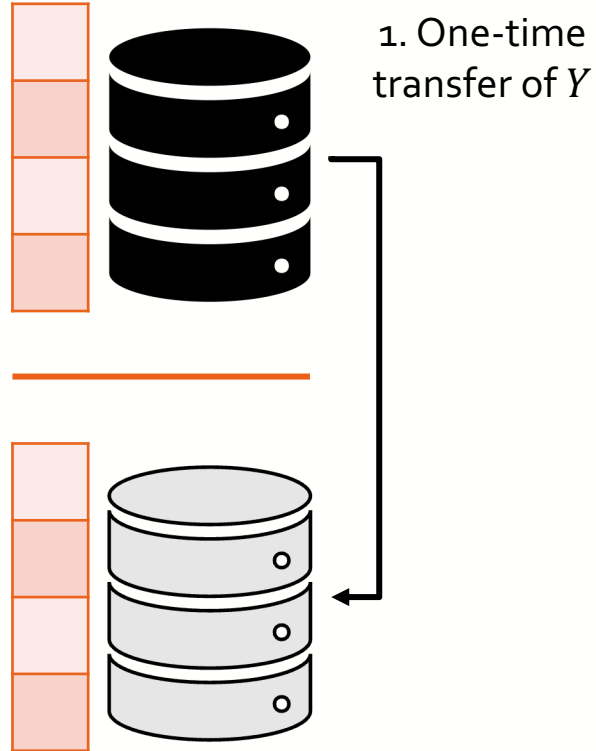
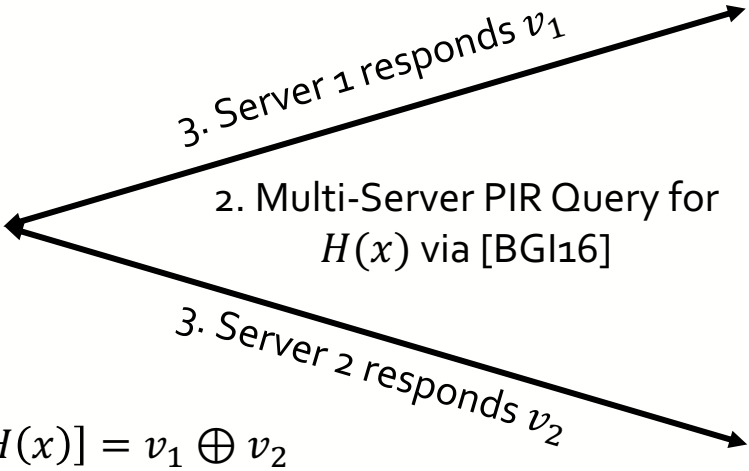
single-server PIR

2/multi-server PIR

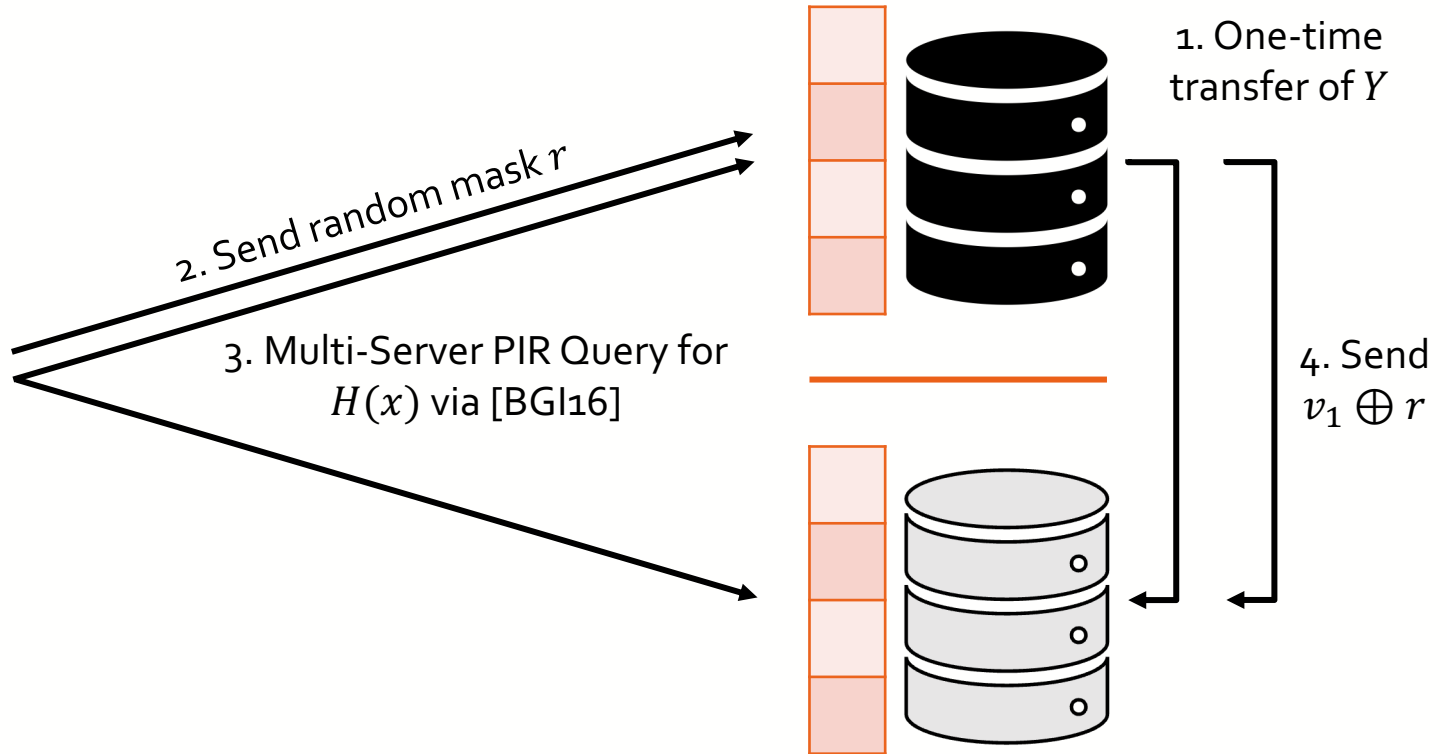
# "PIR-PSI" [DRRT18] – General Idea



Issue: no server input privacy

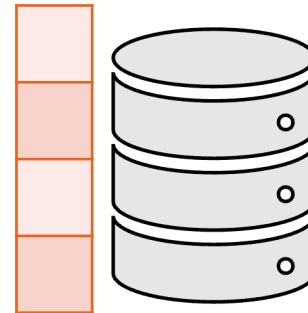
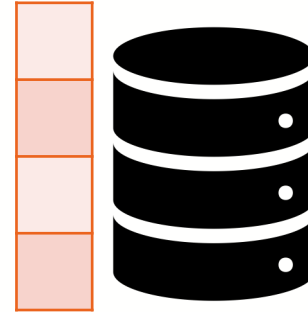
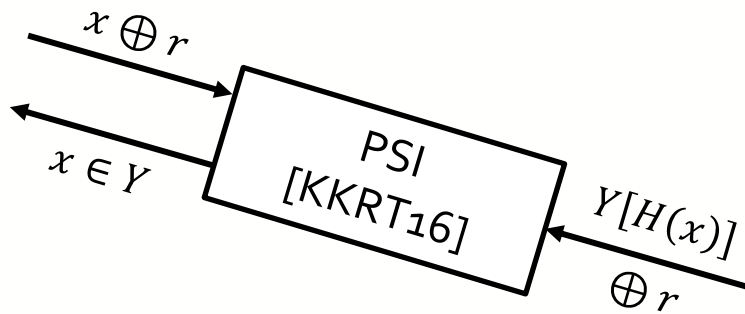


# "PIR-PSI" [DRRT18] – PIR Step





# "PIR-PSI" [DRRT18] – PSI Step

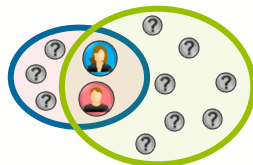


Server 2 knows  
masked PIR result  
 $v_1 \oplus v_2 \oplus r =$   
 $Y[H(x)] \oplus r$

# OPRF-based uPSI w/ PIR Lookup [HSW23]



Client inputs  $x_{i \in \{1, \dots, n\}}$



Instantiated in [HSW23] with  
"Offline/Online"-PIR [KC21]



## 1.1 Server Setup Phase

Client-independent Precomputation Phase to prepare  $PD$

## 1.2 Per-Client Setup Phase

Client-specific Precomputation Phase,  $O(\sqrt{|Y|})$

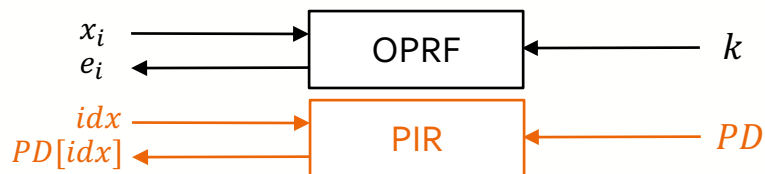
## 2. Base Phase

Client-specific OPRF Setup Phase,  $O(|X|)$

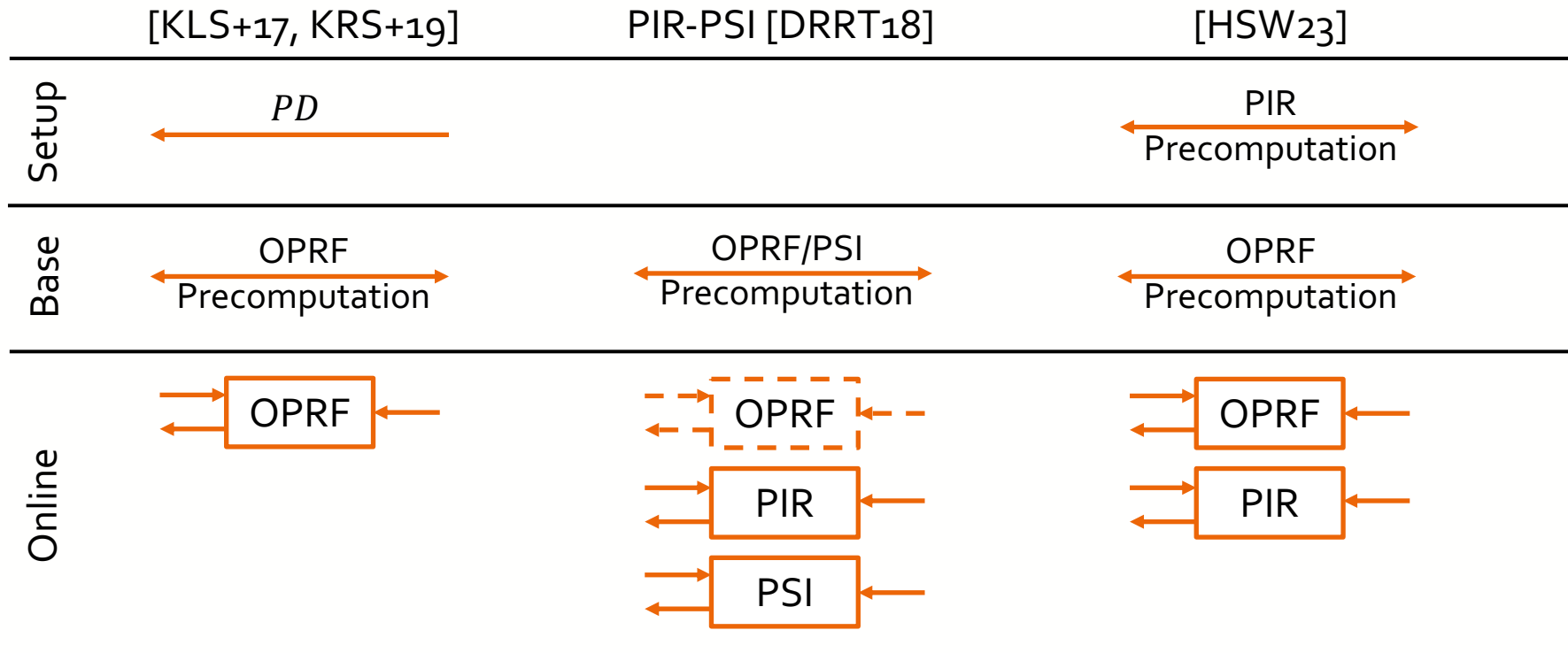
## 3. Online Phase ( $O(|X| + \log(|Y|))$ )

Compute  $idx$  from  $e_i$

Check if  $e_i$  is in  $PD$



# Comparison of OPRF-based $\mu$ PSI – Conceptual



# Comparison of OPRF-based $\mu$ PSI – Performance



Protocol	$ Y $	Time [s]			Comm. [MiB]	
		Server Setup (per client)	Client Setup	Online	Setup	Online
[KRS+19]	$2^{28}$	-	15.17	0.63	1072.14	2.06
[DRRT18]		-	-	13.22	-	5.05
[HSW23]		63.71	35.26	1.08	66.00	4.02
	$2^{31}$	525.09	286.78	1.37	264.00	4.72

Disclaimer: Some results are “cherry picked”; for client sets with size  $|X| = 2^{10}$ ; gentle multi-threading optimizations (4 client + 8 server threads) are considered for [DRRT18, HSW23]

# OPRF-based uPSI w/ Single-Server PIR?

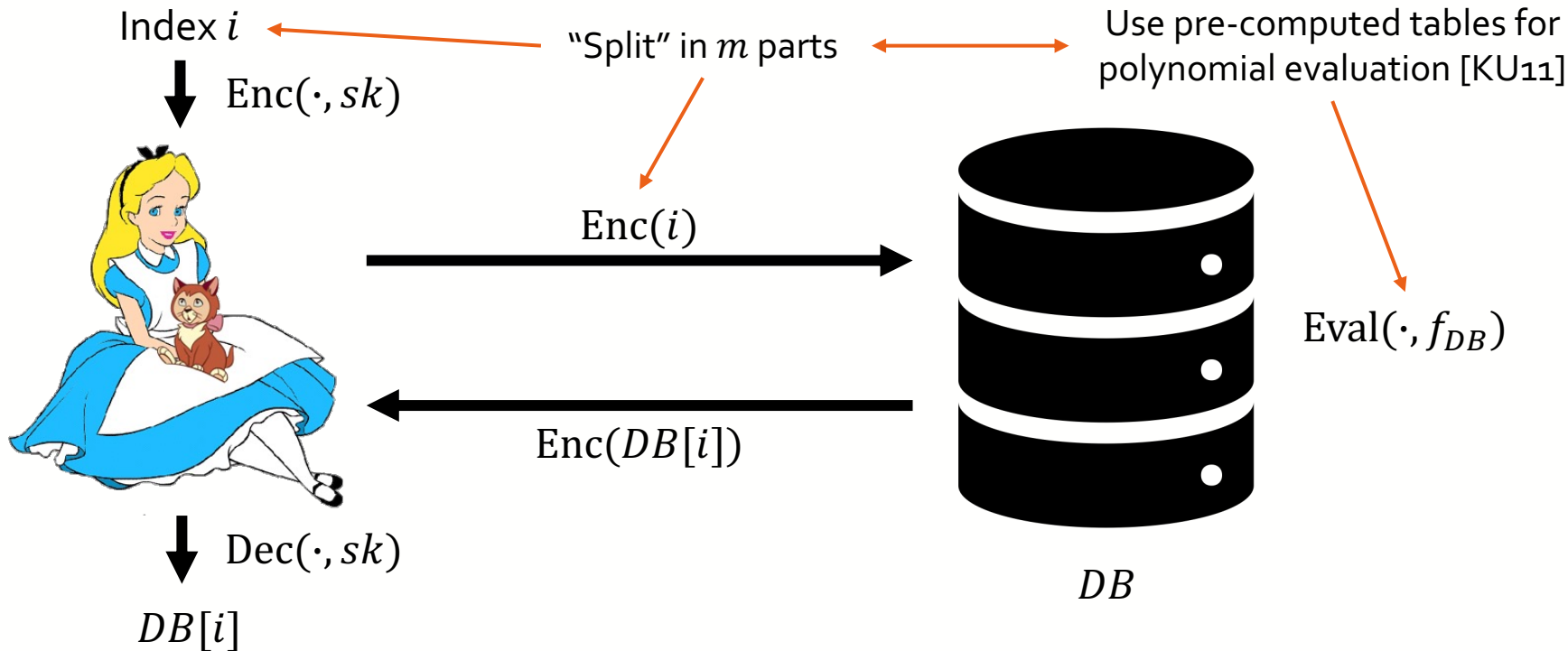


Type	Offline (client-ind.)		Offline (client-dependent)			Online		SOTA Example
	Comp	Storage	Comp	Comm	Hint size	Comp	Comm	
Stateless PIR	$\tilde{O}(N)$	$\tilde{O}(N)$	–			$\tilde{O}(N)$	$\tilde{O}(1)$	[MCR21]
Stateful PIR	$\tilde{O}(N)$	$\tilde{O}(N)$	$\tilde{O}(N)$	$\tilde{O}(\sqrt{N})$	$\tilde{O}(\sqrt{N})$	$\tilde{O}(\sqrt{N})$	$\tilde{O}(1)$	[ZLTS23]
DEPIR	$\tilde{O}(N)$	$\tilde{O}(N)$	–			$\tilde{O}(1)$	$\tilde{O}(1)$	[LMW23]

[LLMT24]: DH-style OPRF [JL10] + SimplePIR [HHC+23]



# Doubly-Efficient PIR (DEPIR) [LMW23]



# Doubly-Efficient PIR (DEPIR) – Concrete Performance



		Database Size			
		46,376	142,506	15,020,334	185,250,786
[OPPW <sub>24a</sub> ]	Total Storage [TB]	0.02	2.04	873.98	411,745.17
	Total Queries [ $\times 2^{30}$ ]	73.2	45	3154	3175
	Run-Time [min]	8	459	(impossible to benchmark)	
[OPPW <sub>24b</sub> ]	Total Storage [TB]	0.04	2.93	750.92	394,474.94
	Total Queries [ $\times 2^{30}$ ]	7	5	134	335
	Run-Time [min]	18	104	(impossible to benchmark)	

# Summary



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# Conclusion



uPSI  
Applications +  
Constructions



uPSI  
Performance



# Thank You!

Get in touch: [christian.weinert@rhul.ac.uk](mailto:christian.weinert@rhul.ac.uk)



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# Literature (1/3)



- [App21] Apple Inc.: "CSAM Detection - Technical Summary". URL: [https://www.apple.com/child-safety/pdf/CSAM\\_Detection\\_Technical\\_Summary.pdf](https://www.apple.com/child-safety/pdf/CSAM_Detection_Technical_Summary.pdf)
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- [DRRT18] D. Demmler, P. Rindal, M. Rosulek, N. Trieu: "PIR-PSI: Scaling Private Contact Discovery". In *PoPETS'18*.
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- [HSW23] L. Hetz, T. Schneider, C. Weinert: "Scaling Mobile Private Contact Discovery to Billions of Users". In *ESORICS'23*.
- [JL10] S. Jarecki, X. Liu: "Fast Secure Computation of Set Intersection". In *SCN'10*.
- [KC21] D. Kogan, H. Corrigan-Gibbs: "Private Blocklist Lookups with Checklist". In *USENIX Sec'21*.

# Literature (2/3)



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# London Crypto Day 2024



ROYAL  
HOLLOWAY  
UNIVERSITY  
OF LONDON

**Date:** Friday 15th November 2024

**Location:** Enigma Room, The Alan Turing Institute (British Library)

## Confirmed Speakers:

- Andrew Mendelsohn, *Imperial College*
- Dave Buckley, *OpenMined*
- Eamonn Postlethwaite, *KCL*
- Maria Corte-Real Santos, *UCL*
- Lydia Garms, *EY*
- Saqib Kakvi, *RHUL*



<https://sites.google.com/view/london-crypto-day/>